

Quick sort

x — x — x

↳ Quick sort is also an important sorting technique used to sort the given data.

↳ Quick sort technique based on divide and conquer strategy

Divide: Partition the array $A[i--j]$ into subarrays $A[i--q-1]$ and $A[q+1--j]$ such that each element of $A[i--q-1] < A[q]$ and each element of $A[q+1--j] > A[q]$, where $A[q]$ is the pivot element. [Generally we take first or last element of the array as a pivot element].

conquer: Sort the two subarrays $A[i--q-1]$ and $A[q+1--j]$ by recursive call to quicksort.

combine: Since the subarrays are sorted in place, no work is needed to combine them: the entire array $A[i--j]$ is now sorted.

↳ $(\log_2 n)$ Pass are required to sort n elements

↳ Best case Time complexity: $O(n \log_2 n)$

↳ worst case Time complexity: $O(n^2)$

↳ Average case Time complexity: $O(n \log_2 n)$.

Rough Idea.

(2)

I/P:

8	10	6	12	7	16	11
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Let's take first element of the array as the pivot element

Pass 1:

7	6	8	10	12	16	11
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~~sorted~~

Pass 2:

6	7	8	10	12	16	11
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 element 7, 8 and 10 sorted

Pass 3: 6 7 8

10	12	16
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 - sorted

Algorithm:

Quick sort

quicksort (A, P, r) , P is a initial index of Array and r is a last index of Array.

```

{
  if (P < r)
  {
    q = partition(A, P, r);
    quicksort(A, P, q-1);
    quicksort(A, q+1, r);
  }
}

```

partition (A, P, r)

```

{
  x = A[r] (pivot element)
  i = P-1
  for (j = P to r-1)
  {
    if (A[j] <= x)
    {
      i = i+1
    }
  }
  swap A[i] with A[j]
  swap A[i+1] with A[r]
  return i+1
}

```